

Modal Logics of Cayley Graphs

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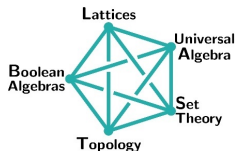
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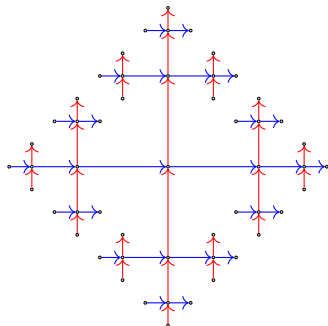
Cayley graphs

Definition

Suppose that G is a group and S is a generating set of G . The *Cayley graph* $\Gamma = \Gamma(G, S)$ is a colored directed graph constructed as follows:

- 1 Each element g of G corresponds to a vertex: the vertex set $V(\Gamma) = G$.
- 2 Each generator s of S is assigned a color c_s .
- 3 For any $g \in G$, $s \in S$ the vertices corresponding to the elements g and gs are joined by a directed edge of color c_s .

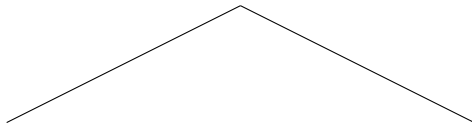
Example



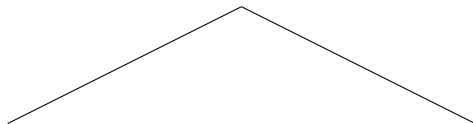
$$F = \langle x, y \mid \emptyset \rangle$$

Cayley graphs as frames

Cayley graph $\Gamma(G, S)$, $|S| = n$

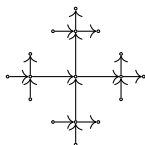


Cayley graphs as frames

Cayley graph $\Gamma(G, S)$, $|S| = n$ Simple Cayley Frame: CayleySimple_n

$$F = (G, \mathcal{R})$$

One relation for all colors

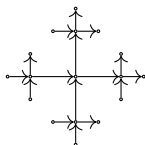
One modality \square

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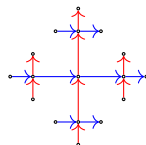
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One relation for all colors

One modality \square Multimodal Cayley Frame: Cayley_S

$$F = (G, \{\mathcal{R}_s\}_{s \in S})$$

Each color defines a relation

Modalities \square_s for each $s \in S$

Four questions

	Simple	Multimodal
General	Given $n \in \mathbb{N}$, $\text{Log}(\text{CayleySimple}_n) = ?$	Given S , $\text{Log}(\text{Cayley}_S) = ?$
Restricted	Given $G = \langle S \mid R \rangle$, $F = (G, \mathcal{R})$, $\text{Log}(F) = ?$	Given $G = \langle S \mid R \rangle$, $F = (G, \{\mathcal{R}_s\}_{s \in S})$, $\text{Log}(F) = ?$

Logic of all multimodal Cayley frames

Definition

The modal logic SL_S is $K_S + \{\Box_s p \leftrightarrow \Diamond_s p\}_{s \in S}$.

Proposition

$Cayley_S \models SL_S$.

Logic of all multimodal Cayley frames

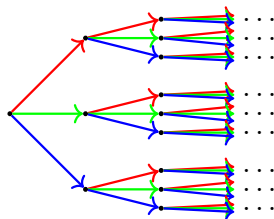
Definition

The modal logic SL_S is $K_S + \{\Box_s p \leftrightarrow \Diamond_s p\}_{s \in S}$.

Proposition

$Cayley_S \models SL_S$.

Each element $g \in G$ has exactly one \mathcal{R}_s -arrow (g, gs) for any $s \in S$.



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$\text{Log}(\text{Cayleys}) \subseteq SL_S.$

Logic of all multimodal Cayley frames

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$$\text{Log}(\text{Cayley}_S) \subseteq SL_S.$$

The logic SL_S is canonical. Consider its canonical frame. The unraveling of any of its point-generated subframes embeds into the Cayley graph of a free group as a subframe.

$$\text{Cayley graph of } \langle S \mid \emptyset \rangle \xrightarrow{\uparrow e} F \uparrow e \longleftrightarrow F_{SL_S}^\# \longrightarrow F_{SL_S}$$

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Corollary

$$\text{Log}(\text{Cayley}_S) = \text{Log}(\text{Multimodal Cayley Frame of } \langle S \mid \emptyset \rangle) = SL_S.$$

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Consider a group $G = \langle S \mid R \rangle$ and its multimodal Cayley frame $F = (W, \{\mathcal{R}_s\})$.

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Idea: translate relations R into modal formulas.

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Seems trivial...

$$xy = yx$$

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$$xyz = e$$

$$\Diamond_x \Diamond_y \Diamond_z p \leftrightarrow p$$

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???

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$xx = y$	
$xy^{-1}xy^{-1} = e$???

Positive relations

Definition

- 1 A word $w = s_1^{\varepsilon_1} s_2^{\varepsilon_2} \dots s_n^{\varepsilon_n}$ in S is positive if $\forall i \in \{1, \dots, n\} \varepsilon_i = 1$, i.e. no inverses of elements of S occur in w .
- 2 For any words w_1, w_2 in S , we call the equation $w_1 = w_2$ positive, if both w_1 and w_2 are positive.
- 3 Let $\langle S \mid R \rangle$ be a group presentation. We call a relation $r \in R$ positive, if r is equivalent to a positive equation.

Examples

- $xy = yx$, $xyz = e$, $xy^{-1}x = e$ are positive;
- $xy^{-1}xy^{-1} = e$ is not positive.

Translation

A translation for positive equation is defined as follows:

- ① $e^\#$ is an empty word;
- ② $s^\# := \diamond_s$ for any $s \in S$;
- ③ $(As)^\# := A^\#s^\#$ for any positive word A in S and any $s \in S$;
- ④ $(A = B)^\# := (A^\#p \leftrightarrow B^\#p)$ for any positive words A, B in S .

Proposition

If G is a group and a positive equation r is true in G , then $r^\#$ is valid in the multimodal Cayley frame of G .

Corollary

If $G = \langle S \mid R \rangle$ and $S = S^{-1}$, then $\{r^\# \mid r \in R\}$ is valid in the multimodal Cayley frame of G .

Proposition

If $G = \langle S \mid R \rangle$, R contains only positive relations, and F the multimodal Cayley frame of G . Then

$$\text{Log}(F) = SL_S + \{r^\# \mid r \in R\}.$$

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For any $r \in R$, $r^\#$ is a Sahlqvist formula, therefore this logic is canonical. We construct a bounded morphism from F to the rooted canonical frame $F_c \uparrow w_0$.

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- ① $f(e) = w_0$;
- ② If $f(g)$ is already set, let $f(gs)$ be the unique element in $\mathcal{R}_s(f(g))$.

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If some $g \in G$ is discovered twice, then

$$g = s_1 s_2 \dots s_m = s'_1 s'_2 \dots s'_k.$$

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But at the same time

$$F_c \models \diamond_{s_1} \diamond_{s_2} \dots r_{s_m} p \leftrightarrow \diamond_{s'_1} \diamond_{s'_2} \dots \diamond_{s'_k} p.$$

Non-positive relations

If r is not positive, it does not affect $\text{Log}(F)$.

Example

$$G = \langle x, y \mid xy^{-1}xy^{-1} = e \rangle; \quad F = (G, \{\mathcal{R}_x, \mathcal{R}_y\}).$$

$$\text{Log}(F) \subseteq \text{Log}(F \uparrow e)$$

Consider the rooted frame $F \uparrow e$. Its points are

$$g = xx \dots xy y \dots yxx \dots xy y \dots y \dots$$

No pattern $xy^{-1}xy^{-1}$ occurs.

Non-positive relations

Definition

Let S be a generating set of a group, and let R be a set of relations in S . We denote $Pos(R)$ the minimal set R' of positive relations that satisfies the following conditions:

- 1 Any positive equality that is true in $\langle S \mid R \rangle$ is also true in $\langle S \mid R' \rangle$.
- 2 Any equality that is true in $\langle S \mid R' \rangle$ is also true in $\langle S \mid R \rangle$.

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Proposition

Let $G = \langle S \mid R \rangle$, where S is finite. Let F be the multimodal Cayley frame of $\Gamma(G, S)$.

Consider the group $G' = \langle S \mid Pos(R) \rangle$, and the multimodal Cayley frame F' of $\Gamma(G', S)$.

The point-generated subframes $F \uparrow e$ and $F' \uparrow e$ are isomorphic.

Main result for multimodal frames

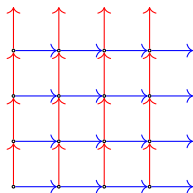
Theorem

Let $G = \langle S \mid R \rangle$, where S is finite. Let $F = (G, \{\mathcal{R}_s\}_{s \in S})$ be the multimodal Cayley frame for G . Then

$$\text{Log } F = KC_S + \{r^\# \mid r \in \text{Pos}(R)\}.$$

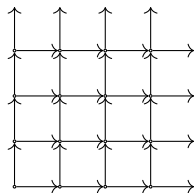
Simple Cayley frames. The difference

$$G = \langle x, y \mid xy = yx \rangle$$



Multimodal Cayley frame

$$F_M = (G, \{\mathcal{R}_x, \mathcal{R}_y\})$$



Simple Cayley frame

$$F_S = (G, \mathcal{R})$$

Frame conditions for simple Cayley frames

$$G = \langle S \mid R \rangle, |S| = n, F = (G, \mathcal{R})$$

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- Serial: $AD = \diamond T$
- Bounded branching:

$$Alt_n = \square \left(\bigvee_{i=1}^{n+1} p_i \right) \rightarrow \diamond \left(\bigvee_{i=1}^{n+1} \bigvee_{j=i+1}^{n+1} (p_i \wedge p_j) \right)$$

Logic of all simple Cayley frames

Theorem

$$\begin{aligned}\text{Log}(\text{CayleySimple}_n) &= \text{Log}(\text{Simple Cayley Frame for } \langle S_n \mid \emptyset \rangle) = \\ &= K + AD + \text{Alt}_n.\end{aligned}$$

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\supseteq : Frame conditions

\subseteq : Canonical frame and truth-preserving operations

$$\text{Cayley graph of } \langle S_n \mid \emptyset \rangle \xrightarrow{\uparrow e} F \uparrow e \longrightarrow F_{K+AD+\text{Alt}_n}^\# \longrightarrow F_{K+AD+\text{Alt}_n}$$

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Can we again translate relations R into modal formulas?

Logic of a simple Cayley frame

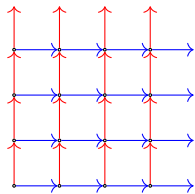
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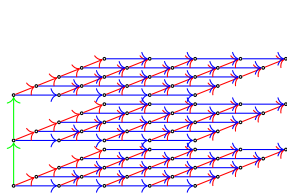
No.

$$G_1 = \langle x, y \mid xy = yx \rangle$$



$$F_{G_1} \models \Diamond \Box p \rightarrow \Box \Diamond p$$

$$G_2 = \langle x, y, z \mid xy = yx \rangle$$



$$F_{G_2} \not\models \Diamond \Box p \rightarrow \Box \Diamond p$$

Negative result on simple Cayley frames

Theorem

There exist no modal logic L and translation function τ that maps group relations to modal formulas, such that for any group $G = \langle S \mid R \rangle$ and its simple Cayley frame $F = (G, \mathcal{R})$,

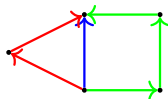
$$\text{Log}(F) = L + \bigwedge_{r \in R} \tau(r).$$

The same is true even if we replace R with $\text{Pos}(R)$.

A family of validities in a simple Cayley frame

Example

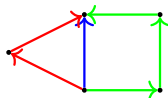
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Frame condition: $\forall x \quad |\mathcal{R}(x) \cap \mathcal{R}^2(x) \cap \mathcal{R}^3(x)| = 1$

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- ③ Let \sim_I be an equivalence relation on G :

$$g_1 \sim_I g_2 \iff L(PW(g_1)) = L(PW(g_2)).$$

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- ③ Let \sim_I be an equivalence relation on G :

$$g_1 \sim_I g_2 \iff L(PW(g_1)) = L(PW(g_2)).$$

- ④ For each $C \in G / \sim_I$, denote $L(C) = L(PW(g))$ for any $g \in C$.

A family of validities in a simple Cayley frame

Consider any class of equivalence $C \in G / \sim_I$.

- If C is infinite, we may define

$$\forall x \left| \bigcap_{I \in L(C)} \mathcal{R}'(x) \right| \geq 1;$$

- If C is finite, we may define

$$\forall x \ 1 \leq \left| \bigcap_{I \in L(C)} \mathcal{R}'(x) \right| \leq |C|;$$

A family of validities in a simple Cayley frame

For an infinite $C \in G / \sim_I$ such that $L(C) = \{l_1, \dots, l_n\}$ is finite, define

$$\phi^\infty(C) = \bigwedge_{i=1}^n \square^{l_i} p_i \rightarrow \diamond^{l_1} \bigwedge_{i=1}^n p_i.$$

For a finite $C \in G / \sim_I$ such that $L(C) = \{l_1, \dots, l_n\}$ is finite, define

$$\begin{aligned} \phi(C) &= \left(\left(\bigwedge_{i=1}^n \square^{l_i} p_i \right) \wedge \left(\bigwedge_{i=1}^n \left(p_i \rightarrow \bigvee_{j=1}^{|C|+1} q_j \right) \right) \right) \rightarrow \\ &\rightarrow \diamond^{l_1} \left(\left(\bigwedge_{i=1}^n p_i \right) \wedge \left(\bigvee_{j=1}^{|C|+1} \bigvee_{j'=j+1}^{|C|+1} q_j \wedge q_{j'} \right) \right). \end{aligned}$$

A family of validities in a simple Cayley frame

Theorem

Let $F = (G, \mathcal{R})$ be the simple Cayley frame for a group $G = \langle S \mid R \rangle$. Then

- ① For any finite $C \in G / \overset{!}{\sim}$ such that $L(C)$ is finite,

$$F \models \phi(C)$$

- ② For any infinite $C \in G / \overset{!}{\sim}$ such that $L(C)$ is finite,

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- ② For any infinite $C \in G / \sim^I$ such that $L(C)$ is finite,

$$F \models \phi^\infty(C)$$

If $L(C)$ is infinite, formulas $\phi(C)$ where $L(C)$ is replaced with any finite subset $L' \subseteq L(C)$ are valid in F .

Results

	Simple	Multimodal
General	$K + AD + Alt_n$	SL_S
Restricted	$\{\phi(C) \mid C \in G / \sim^I \text{ (finite)}\}$ $\{\phi^\infty(C) \mid C \in G / \sim^I \text{ (infinite)}\}$...	$SL_S + \{r^\# \mid r \in Pos(R)\}$

Thank you for your attention!