Lattice Theory Lecture 2

Distributive lattices

John Harding

New Mexico State University www.math.nmsu.edu/~JohnHarding.html

jharding@nmsu.edu

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Distributive lattices

Distributive law	for all x, y, z	$x \lor (y \land z) = (x \lor y) \land (x \lor z)$
Modular law	if $x \leq z$ then	$x \lor (y \land z) = (x \lor y) \land (x \lor z)$

Definition The lattices M_5 and N_5 are as follows:



Note M_5 is **M**odular, not distributive, and N_5 is **N**on-modular. Both have 5 elements.

Recognizing distributive lattices

Theorem Let L be a lattice.

- 1. L is modular iff N_5 is not a sublattice of L
- 2. L is distributive iff neither M_5, N_5 is a sublattice of L

Proof The " \Rightarrow " direction of each is obvious. For 1 " \Leftarrow " if *L* is not modular, there are x < z with $x \lor (y \land z) < (x \lor y) \land (x \lor z)$ (why?) Then the following is a sublattice of *L*.



Exercise

Give the details that the figure on the previous page is a sublattice.

Do the 2 " \Leftarrow " direction.

The lattice N_5 is "projective" in lattices, meaning that if L is a lattice and $f: L \rightarrow N_5$ is an onto lattice homomorphism, then there is a one-one lattice homomorphism $g: N_5 \rightarrow L$ with $f \circ g = id$.

Complements

Definition Elements x, y of a bounded lattice L are complements if $x \wedge y = 0$ and $x \vee y = 1$.

In general, an element might have no complements, or many.



Complements

Theorem In a bounded distributive lattice, an element has at most one complement.

Pf Suppose y, z are complements of x. Then

$$y = y \land (x \lor z) = (y \land x) \lor (y \land z) = y \land z$$
$$z = z \land (x \lor y) = (z \land x) \lor (z \land y) = y \land z$$

Surprisingly, a finite lattice where each element has exactly one complement is distributive! But not so in the infinite case.

Boolean algebras

Definition A Boolean algebra $(B, \land, \lor, ', 0, 1)$ is an algebra of type 2,2,1,0,0 where

- 1. B is abounded distributive lattice
- 2. x' is a complement of x for each $x \in B$

Note The difference between a complemented distributive lattice and a Boolean algebra is what we consider to be a subalgebra. A subalgebra of a Boolean algebra must include complements.

Properties of Boolean algebras

Proposition In any Boolean algebra

1.
$$(x \land y)' = x' \lor y'$$

2. $(x \lor y)' = x' \land y'$

3.
$$x'' = x$$

Note These are called De Morgan's laws.

Exercise Prove these. For (1) show $x \wedge y$ and $x' \vee y'$ are complements.

Complements

Definition For L a lattice and $a, b \in L$ with $a \leq b$ the interval [a, b] is the sublattice of L given by

$$[a,b] = \{x : a \le x \le b\}$$

Proposition Each interval [a, b] in a complemented distributive lattice *L* is complemented with the complement of *x* being the element $x^{\#}$ given by

$$x^{\#} = (x' \wedge b) \vee a$$

We say that L is relatively complemented when its intervals are complemented. The above result works for modular lattices too.

Complements

Definition For a bounded distributive lattice L, let its center C(L) be the set of all complemented elements of L.

Proposition The center of L is a sublattice of L.

Proposition $c \in C(L)$ gives an isomorphism $\varphi : L \rightarrow [0, c] \times [0, c']$ where

$$\varphi(x) = (x \wedge c, x \wedge c')$$

Further, each direct product decomposition of L arises this way.

Pf Define $\psi : [0, c] \times [0, c'] \rightarrow L$ by $\psi(p, q) = p \lor q$ and show it is inverse to ϕ . For the further comment, if $L = A \times B$, then c = (1, 0) is in C(L).

Definition An ideal of a lattice L is a subset $I \subseteq L$ where

- 1. if $y \in I$ and $x \leq y$, then $x \in I$
- 2. if $x, y \in I$ then $x \lor y \in I$

Definition A filter of a lattice L is a subset $F \subseteq L$ where

1. if
$$x \in F$$
 and $x \leq y$, then $y \in F$

2. if $x, y \in F$ then $x \wedge y \in F$

Definition Let $\mathcal{I}(L)$ be the set of ideals of L partially ordered by \subseteq .

Definition For any L its ideal lattice $\mathcal{I}(L)$ is a complete lattice with meets given by intersections. The join of two ideals I and J is

 $I \lor J = \{x : x \le a \lor b \text{ for some } a \in I, b \in J\}$

Pf We check that the intersection of ideals is an ideal. Then the ideals are a closure system, hence a complete lattice. We check that the description above is an ideal, and then must be the smallest ideal containing I, J.

Exercise Do the details.

Definition For a lattice *L* and $a \in L$, the principal ideal and principal filter generated by *a* are $\downarrow a = \{x : x \le a\}$ and $\uparrow a = \{x : a \le x\}$.

Proposition $\phi(a) = \downarrow a$ is a lattice embedding of L into $\mathcal{I}(L)$.

Pf Clearly
$$\downarrow a \cap \downarrow b = \downarrow (a \land b)$$
 and $\uparrow a \lor \uparrow b = \uparrow (a \lor b)$

Definition An ideal of L is trivial if it is either empty or all of L.

Proposition If *L* is distributive, so is $\mathcal{I}(L)$.

Pf Let I, J, K be ideals. Always $I \lor (J \land K) \subseteq (I \lor J) \land (I \lor K)$.

- If $x \in RHS$, then
- $x \in I \lor J$ and $x \in I \lor K$
- exist $a_1, a_2 \in I$, $b \in J$, $c \in K$ with $x \le a_1 \lor b$ and $x \le a_2 \lor c$

•
$$a = a_1 \lor a_2 \in I$$

- $x \leq (a \lor b) \land (a \lor c) = a \lor (b \land c)$
- *x* ∈ LHS.

A fancier version of this shows ...

Theorem For any lattice L, the ideal lattice $\mathcal{I}(L)$ satisfies exactly the same equations as L.

There are several results constructing $\mathcal{I}(L)$ as a homomorphic image of a subalgebra of an ultrapower of *L*.

Note These results are for lattice equations. For a Boolean algebra B, its ideal lattice $\mathcal{I}(B)$ need not be complemented.

For certain lattices, ideals play a role similar to that of normal subgroups for groups.

Proposition For I an ideal of a distributive lattice L, there is a congruence θ_I of L where

$$\theta_I = \{(a, b) : a \lor x = b \lor x \text{ for some } x \in I\}$$

If L is sectionally complemented, these are all of its congruences.

Exercise Prove the first statement using the definition of a congruence. The second is a bit harder.

Prime ideals

Definition An ideal P of a lattice L is prime if for any $a, b \in L$

$$a \land b \in P \implies a \in P \text{ or } b \in P$$

Proposition An ideal P is prime iff $L \setminus P$ is a filter.



A prime ideal and its complementary filter split the lattice in two.

Prime ideals

Definition Let 2 be the 2-element lattice

Proposition For *P* a prime ideal of a distributive lattice *D*, there is a homomorphism $\varphi_p: D \to 2$ where

$$\varphi_p(x) = \begin{cases} 0 & \text{if } x \in P \\ 1 & \text{if } x \notin P \end{cases}$$



The Prime Ideal Theorem

Theorem Let *I* be an ideal of a distributive lattice *D*, and *F* be a filter of *D* with $I \cap F = \emptyset$. Then there is a prime ideal *P* with $I \subseteq P$ and *F* disjoint from *P*.



The Prime Ideal Theorem

Pf Let $\mathfrak{X} = \{J : I \subseteq J \in \mathcal{I}(L) \text{ and } J \cap F = \emptyset\}.$

- \mathfrak{X} is non-empty and closed under unions of chains
- By Zorn's Lemma $\mathfrak X$ has a maximal member P
- *P* is an ideal, $I \subseteq P$ and $P \cap F = \emptyset$
- Suppose a, b ∉ P
- by maximality $\downarrow a \lor P, \downarrow b \lor P \notin \mathfrak{X}$
- exist $x_1, x_2 \in P$ with $a \lor x_1 \in F$ and $b \lor x_2 \in F$
- $x = x_1 \lor x_2 \in P$
- $(a \lor x) \land (b \lor x) = (a \land b) \lor x \in F$
- since $P \cap F = \emptyset$ then $a \wedge b \notin P$.

The Prime Ideal Theorem

If we consider the case of propositional statements ...

An ideal of statements is a collection that we can sensibly decide to assign value FALSE. A filter is a collection we can sensibly assign the value TRUE.

The prime ideal theorem says that we can do both together while assigning T_{RUE} or F_{ALSE} to every proposition in a consistent way! This is a bit remarkable if you think of it.

Definition For a distributive lattice D, let $\beta(D)$ be the set of all non-trivial prime ideals of D.

Definition For D a distributive lattice and $a \in D$ set

$$\beta(a) = \{P \in \beta(D) : a \notin P\}$$

Proposition For D a distributive lattice and $a, b \in D$

1.
$$\beta(a \land b) = \beta(a) \cap \beta(b)$$

2. $\beta(a \lor b) = \beta(a) \cup \beta(b)$

Exercise Prove this proposition.

Theorem Any distributive lattice D is isomorphic to a sublattice of the power set $\mathcal{P}(X)$ of the set $X = \beta(D)$.

Pf The map $\beta: D \to \mathcal{P}(X)$ preserves \land and \lor . It remains to show it is one-one.

- Let *a* ≠ *b*
- Either $a \notin b$ or $b \notin a$
- Assume b ≰ a
- Then ↓a and ↑b are a disjoint ideal and filter
- There is a prime ideal P with $\downarrow a \subseteq P$ and $\uparrow b \cap P = \emptyset$
- a ∈ P and b ∉ P
- $\beta(a) \neq \beta(b)$

Proposition For any set X, the power set $\mathcal{P}(X)$ is isomorphic to the power 2^X , all functions from X to 2 with pointwise operations.

Pf Let χ be the map from $\mathcal{P}(X)$ to 2^X sending a subset $A \subseteq X$ to its characteristic function χ_A where

$$\chi_A(x) = \begin{cases} 0 & \text{if } x \notin A \\ 1 & \text{if } 1 \in A \end{cases}$$

Theorem Any distributive lattice D is isomorphic to a sublattice of the power 2^X where $X = \beta(D)$.

Prime ideals for Boolean algebras

Exercises Let I be a non-trivial ideal of a Boolean algebra B. Show that the following are equivalent.

- 1. *I* is prime
- 2. For each $x \in B$ exactly one of x, x' belongs to *I*.
- 3. *I* is a maximal non-trivial ideal.

Thus $\beta(x') = \beta(B) \setminus \beta(x)$.

Pf Exercise.

For Boolean algebras of equivalence classes of logical statements, prime ideals correspond to consistent assignments of truth and falsehood to the statements.

Prime ideals for Boolean algebras

Corollary Let *B* be a Boolean algebra *B* and $X = \beta(B)$.

- 1. *B* is isomorphic to a subalgebra of the power set $\mathcal{P}(X)$.
- 2. *B* is isomorphic to a subalgebra of the power 2^X .

Remark This is much like Cayley's theorem for groups that says every group G is isomorphic to a subgroup of the group of permutations of a set.

Remark This theorem was proved independently by Stone and Birkhoff in the 1930's. Stone was an analyst and this result plays a key role in functional analysis.

The result that each distributive lattice and each Boolean algebra is a subalgebra of $\mathcal{P}(X)$ has useful consequences.

Example To show that the following holds in each Boolean algebra

$$x \le y \Leftrightarrow x \land y' = 0$$

its enough to verify it for $\mathcal{P}(X)$ where it is $S \subseteq T \Leftrightarrow S \cap T' = \emptyset$.

The result that each is a subalgebra of 2^X is also useful. We switch focus and look at familiar ideas from logic from this perspective.

Definition A term $t(x_1, ..., x_n)$ in the language of lattices is an expression built from \land, \lor and the variables $x_1, ..., x_n$.

Example
$$t(x, y, z) = ((x \land y) \land x) \lor z$$

Definition For a term $t(x_1, ..., x_n)$ and lattice A, the interpretation of t in A is the function $t^A : A^n \to A$ where $t^A(a_1, ..., a_n)$ is the value of the term when its inputs are $a_1, ..., a_n$.

Example $t^A(a, b, c) = ((a \land b) \land a) \lor c$

Truth tables

Definition A truth table is an interpretation of a term in 2.

Example For $t(x, y, z) = ((x \land y) \land x) \lor z$

Decidability

Definition Terms s, t are equivalent if their interpretations s^A, t^A are equal in any distributive lattice A. We then write

$$s(x_1,\ldots,x_n) \equiv t(x_1,\ldots,x_n)$$

Theorem s, t are equivalent iff $s^2 = t^2$.

Pf " \Rightarrow " is vacuous. " \Leftarrow " For A a distributive lattice, $A \le 2^X$. Evaluating s, t in A is componentwise. As s = t in 2, s = t in A.

Remark One says that the equational theory of distributive lattices is decidable. It all works for Boolean algebras too.

Free Algebras

Definition For a set X, a distributive lattice $\mathcal{F}_D(X)$ is called a free distributive lattice over X if

- 1. $\mathcal{F}_D(X)$ is generated by X
- 2. for any distributive lattice D and set mapping $f: X \rightarrow D$



there is a homomorphism $\overline{f} : \mathcal{F}_D(X) \to D$ extending f.

Free Algebras

This is a key notion in algebra, logic, algebraic topology, and computer science. Free groups and free Boolean algebras are defined similarly.

Constructing Free Algebras

Definition Let $\mathcal{T}(X)$ be all terms for distributive lattices whose variables are from X, and let \equiv be the relation of equivalence of terms.

The following theorem from universal algebra holds with obvious modification for groups, rings, lattices, Boolean algebras, and so forth.

Theorem $\mathcal{F}_D(X)$ is equal to $\mathcal{T}(X)/\equiv$.

So $\mathcal{F}_D(X)$, and its Boolean counterpart $\mathcal{F}_B(X)$, are key in logic. They literally are logical propositions modulo logical equivalence.

Free Algebras

For the following, note that the elements of 2^{2^X} are truth tables! Theorem Let X be a finite set.

1. $\mathcal{F}_D(X)$ is isomorphic to a sublattice of 2^{2^X} 2. $\mathcal{F}_B(X)$ is isomorphic to 2^{2^X}

Pf (1) Define
$$\varphi : (\mathcal{T}(X)/\equiv) \longrightarrow 2^{2^X}$$
 by $\varphi(t/\equiv) = t^2$

(2) Since every truth table can be realized by a Boolean algebra term we have that φ is onto in the Boolean case.

Free Algebras

Corollary A subalgebra of a distributive lattice or Boolean algebra generated by n elements is finite.

Pf Such a subalgebra is generated by n elements, and is therefore a homomorphic image of a free algebra on n generators.

Open problem Give a formula for the cardinality of $\mathcal{F}_D(n)$.

Exercise Give an infinite lattice that is generated by 3 elements.

The distributive laws say

$$x \land (y \lor z) = (x \land y) \lor (x \land z)$$
$$x \lor (y \land z) = (x \lor y) \land (x \lor z)$$

Definition A complete lattice D is infinitely meet distributive if it satisfies (1) and infinitely join distributive if it satisfies (2)

1.
$$x \land \bigvee_I y_i = \bigvee_I (x \land y_i)$$

2. $x \lor \land_I y_i = \land_I (x \lor y_i)$

Exercise Show that any finite lattice, any complete chain, and any power set $\mathcal{P}(X)$ satisfies both.

Example The lattice below is complete and distributive, but does not satisfy the infinite meet distributive law.



Then $x \land \forall y_i = x \land 1 = x$ and $\forall (x \land y_i) = 0$.

Proposition The lattice $\mathcal{O}(X)$ of open sets of a topological space X satisfies the infinite meet distributive law but not necessarily the infinite join distributive law.

Pf In $\mathcal{O}(X)$ arbitrary joins are \bigcup and finite meets are \cap , so the result follows from that for sets.

For the failure of infinite join continuity, consider the topological space \mathbb{R} . Set $A = \mathbb{R} \setminus \{0\}$ and $B_n = (-1/n, 1/n)$.

$$A \cup \bigwedge B_n = A \cup \emptyset = A \qquad \bigwedge (A \cup B_n) = \mathbb{R}$$

Here $\bigwedge B_n$ is the interior of its intersection $\{0\}$, which is empty.

There is a stronger version of distributivity involving both infinite joins and meets. To see its nature, lets return to when we were 7.

$$(x_{11} + x_{12}) \cdot (x_{21} + x_{22} + x_{23}) \cdot (x_{31} + x_{32} + x_{33} + x_{34}) = (x_{11} \cdot x_{21} \cdot x_{31}) + (x_{11} \cdot x_{21} \cdot x_{32}) + \dots + (x_{12} \cdot x_{23} \cdot x_{34})$$

There are $24 = 2 \times 3 \times 4$ terms here, one for each choice function.

Definition A complete lattice L is completely distributive if

$$\bigwedge_{I} \bigvee_{J_{i}} x_{ij} = \bigvee_{\alpha \in \prod J_{i}} \bigwedge_{I} x_{i,\alpha(i)}$$

Exercise Show that every finite distributive lattice, every complete chain, and every power set $\mathcal{P}(X)$ is completely distributive.

Exercise Show that the complete distributive law implies the infinite join and meet distributive laws.

Complete Boolean Algebras

Proposition In a complete Boolean algebra B

1.
$$(\forall x_i)' = \land x_i'$$

2. $(\land x_i)' = \lor x_i'$

This works without completeness if we assume one side exists.

Pf (1) $x_i \le y \Leftrightarrow y' \le x'_i$. So if y is the least upper bound of the x_i , then y' is the greatest lower bound of the x'_i . (2) Similar.

Complete Boolean Algebras

Proposition Every complete Boolean algebra satisfies the infinite meet and join distributive laws.

Pf Exercise. Hint: for $\bigvee (x \land y_i) = x \land \bigvee y_i$ trivially LHS \leq RHS. For the other way, it is enough to show that LHS ' \land RHS = 0. Then use LHS' $\leq x' \lor y'_i$ for each *j* to obtain this.

In the final lecture, one more result will be of key importance.

Complete Boolean Algebras

Theorem For a complete Boolean algebra B, these are equivalent.

- 1. B is atomic
- 2. *B* is completely distributive
- 3. *B* is isomorphic to a power set $\mathcal{P}(X)$ for some set *X*.

Pf (Sketch) (1) \Rightarrow (3) \Rightarrow (2) are an exercise. To show (2) \Rightarrow (1) enumerate *B* as x_i and set $x_{i0} = x_i$ and $x_{i1} = x'_i$.

$$x_j < \bigwedge_{I} x_{i\alpha(i)} \Rightarrow \alpha(j) = 1 \Rightarrow x_j \le x'_j \Rightarrow x_j = 0$$

So $\bigwedge_I x_{i\alpha(i)}$ is either 0 or an atom.

$$1 = \bigwedge_{I} \bigvee_{2} x_{ij} = \bigvee_{\alpha \in 2^{I}} \bigwedge_{I} x_{i\alpha(i)}$$

Thanks for listening.

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